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INTRODUCTION TO MONOIDS

Overview

- Monoids (definition, examples)
- Reducers
- Generators
- Benefits of Monoidal Parsing
 - Incremental Parsing (FingerTrees)
 - Parallel Parsing (Associativity)
 - Composing Parsers (Products, Layering)
 - Compressive Parsing (LZ78, Bentley-McIlroy)
- Going Deeper (Seminearrings)

What is a Monoid?

- A Monoid is *any* associative binary operation with a unit.
- Associative: $(a + b) + c = a + (b + c)$
- Unit: $(a + \emptyset) = a = (\emptyset + a)$
- Examples:
 - $((*), 1)$, $((+), \emptyset)$, $(\max, \text{minBound})$,
 $((.), \text{id})$, ...

Monoids as a Typeclass

```
class Monoid m where
```

```
  mempty  :: m
```

```
  mappend :: m -> m -> m
```

```
  mconcat :: [m] -> m
```

```
  mconcat = foldr mappend mempty
```

Built-in Monoids

```
newtype Sum a = Sum a
```

```
instance Num a => Monoid (Sum a) where
```

```
  mempty = Sum 0
```

```
  Sum a `mappend` Sum b = Sum (a + b)
```

```
newtype Endo a = Endo (a -> a)
```

```
instance Monoid (Endo a) where
```

```
  mempty = Endo id
```

```
  Endo f `mappend` Endo g = Endo (f . g)
```

So how can we use them?

- Data.Foldable provides fold and foldMap

```
class Functor t => Foldable t where
```

```
...
```

```
fold :: Monoid m => t m -> m
```

```
foldMap :: Monoid m => (a -> m) -> t a -> m
```

```
fold = foldMap id
```

Monoids are Compositional

```
instance (Monoid m, Monoid n) => Monoid (m,n) where
  mempty = (mempty,mempty)
  (a,b) `mappend` (c,d) = (a `mappend` c, b `mappend` d)
```


But we always pay full price

- Containers are Monoid-oblivious
- Monoids are Container-oblivious

Can we fix that and admit optimized folds?

`(:)` is faster than `(\x xs -> return x ++ xs)`

And what about monotypic containers?

Strict and Lazy ByteStrings, IntSets, etc...

Monoid-specific efficient folds

```
class Monoid m => Reducer c m where
```

```
  unit :: c -> m
```

```
  snoc :: m -> c -> m
```

```
  cons :: c -> m -> m
```

```
c `cons` m = unit c `mappend` m
```

```
m `snoc` c = m `mappend` unit c
```

Simple Reducers

instance Reducer a [a] where

unit a = [a]

cons = (:)

instance Num a => Reducer a (Sum a) where

unit = Sum

instance Reducer (a -> a) (Endo a) where

unit = Endo

Reducers enable faster folds

```
reduceList :: (c `Reducer` m) => [c] -> m  
reduceList = foldr cons mempty
```

```
reduceText :: (Char `Reducer` m) => Text -> m  
reduceText = Text.foldl' snoc mempty
```

Non-Functorial Containers

```
class Generator c where
```

```
  type Elem c :: *
```

```
  mapReduce :: (e `Reducer` m) => (Elem c -> e) -> c -> m
```

```
  ...
```

```
reduce :: (Generator c, Elem c `Reducer` m) => c -> m
```

```
reduce = mapReduce id
```

```
instance Generator [a] where
```

```
  type Elem [a] = a
```

```
  mapReduce f = foldr (cons . f) mempty
```

Container-Specific Folds

```
instance Generator Strict.ByteString where
  type Elem Strict.ByteString = Word8
  mapReduce f = Strict.foldl' (\a b -> snoc a (f b)) mempty
```

```
instance Generator IntSet where
  type Elem IntSet = Int
  mapReduce f = mapReduce f . IntSet.toList
```

```
instance Generator (Set a) where
  type Elem (Set a) = a
  mapReduce f = mapReduce f . Set.toList
```

Parallel ByteString Reduction

instance Generator Lazy.ByteString where

mapReduce f =

 Data.Foldable.fold .

 parMap rwhnf (mapReduce f) .

 Lazy.toChunks

Non-Trivial Monoids/Reducers

- Tracking Accumulated File Position Info
- FingerTree Concatenation
- Delimiting Words
- Parsing UTF8 Bytes into Chars
- Parsing Regular Expressions
- Recognizing Haskell Layout
- Parsing attributed PEG, CFG, and TAGs!

Generator Combinators

```
mapM_ :: (Generator c, Monad m) => (Elem c -> m b) -> c -> m ()
```

```
forM_ :: (Generator c, Monad m) => c -> (Elem c -> m b) -> m ()
```

```
msum :: (Generator c, MonadPlus m, m a ~ Elem c) => c -> m a
```

```
traverse_ :: (Generator c, Applicative f) => (Elem c -> f b) -> c -> f ()
```

```
for_ :: (Generator c, Applicative f) => c -> (Elem c -> f b) -> f ()
```

```
asum :: (Generator c, Alternative f, f a ~ Elem c) => c -> f a
```

```
and :: (Generator c, Elem c ~ Bool) => c -> Bool
```

```
or :: (Generator c, Elem c ~ Bool) => c -> Bool
```

```
any :: Generator c => (Elem c -> Bool) -> c -> Bool
```

```
all :: Generator c => (Elem c -> Bool) -> c -> Bool
```

```
foldMap :: (Monoid m, Generator c) => (Elem c -> m) -> c -> m
```

```
fold :: (Monoid m, Generator c, Elem c ~ m) => c -> m
```

```
toList :: Generator c => c -> [Elem c]
```

```
concatMap :: Generator c => (Elem c -> [b]) -> c -> [b]
```

```
elem :: (Generator c, Eq (Elem c)) => Elem c -> c -> Bool
```

```
filter :: (Generator c, Reducer (Elem c) m) => (Elem c -> Bool) -> c -> m
```

```
filterWith :: (Generator c, Reducer (Elem c) m) => (m -> n) -> (Elem c -> Bool) -> c -> n
```

```
find :: Generator c => (Elem c -> Bool) -> c -> Maybe (Elem c)
```

```
sum :: (Generator c, Num (Elem c)) => c -> Elem c
```

```
product :: (Generator c, Num (Elem c)) => c -> Elem c
```

```
notElem :: (Generator c, Eq (Elem c)) => Elem c -> c -> Bool
```

Generator Combinators

- Most generator combinators just use `mapReduce` or `reduce` on an appropriate monoid.

```
reduceWith f = f . reduce
```

```
mapReduceWith f g = f . mapReduce g
```

```
sum = reduceWith getSum
```

```
and = reduceWith getAll
```

```
any = mapReduceWith getAny
```

```
toList = reduce
```

```
mapM_ = mapReduceWith getAction
```

```
...
```

Example: File Position Delta

- We track the delta of column #s

```
data Delta = Cols Int | ...
```

```
instance Monoid Delta where  
  mempty = Cols 0  
  Cols x `mappend` Cols y = Cols (x + y)
```

```
instance Reducer Delta Char where  
  unit _ = Cols 1
```

-- but what about newlines?

Handling Newlines

- After newline, preceding columns are useless, and we know an absolute column #

```
data Delta = Cols Int | Lines Int Int | ...
```

```
instance Monoid Delta where
```

```
  Lines l _ `mappend` Lines l' c' = Lines (l + l') c'
```

```
  Cols _ `mappend` Lines l' c' = Lines l c'
```

```
  Lines l c `mappend` Cols c' = Lines l (c + c')
```

```
  ...
```

```
instance Reducer Delta where
```

```
  unit '\n' = Lines 1 1
```

```
  unit _ = Cols 1
```

- but what about tabs?

Handling Tabs

```
data Delta = Cols Int | Lines Int Int | Tabs Int Int | ...
```

```
nextTab :: Int -> Int
```

```
nextTab !x = x + (8 - (x - 1) `mod` 8)
```

```
instance Monoid Delta where
```

```
...
```

```
Lines l c `mappend` Tab x y = Lines l (nextTab (c + x) + y)
```

```
Tab {} `mappend` l@Lines {} = l
```

```
Cols x `mappend` Tab x' y = Tab (x + x') y
```

```
Tab x y `mappend` Cols y' = Tab x (y + y')
```

```
Tab x y `mappend` Tab x' y' = Tab x (nextTab (y + x') + y')
```

```
instance Reducer Char Delta where
```

```
unit '\t' = Tab 0 0
```

```
unit '\n' = Line 1 1
```

```
unit _ = Cols 1
```

#line Directives

```
data Delta =  
  = Pos !ByteString !Int !Int  
  | Line !Int !Int  
  | Col !Int  
  | Tab !Int !Int
```

Delta

instance Monoid Delta where

 mempty = Cols 0

 Cols c `mappend` Cols d = Cols (c + d)

 Cols c `mappend` Tab x y = Tab (c + x) y

 Lines l c `mappend` Cols d = Lines l (c + d)

 Lines l _ `mappend` Lines m d = Lines (l + m) d

 Lines l c `mappend` Tab x y = Lines l (nextTab (c + x) + y)

 Tab x y `mappend` Cols d = Tab x (y + d)

 Tab x y `mappend` Tab x' y' = Tab x (nextTab (y + x') + y')

 Pos f l _ `mappend` Lines m d = Pos f (l + m) d

 Pos f l c `mappend` Cols d = Pos f l (c + d)

 Pos f l c `mappend` Tab x y = Pos f l (nextTab (c + x) + y)

 _ `mappend` other = other

data Delta

 = Pos S.ByteString !Int !Int

 | Lines !Int !Int

 | Tab !Int !Int

 | Cols !Int

 deriving (Eq, Show, Data, Typeable)

nextTab :: Int -> Int

nextTab x = x + (8 - x `mod` 8)

instance Reducer Char Delta where

 unit '\n' = Lines 1 1

 unit '\t' = Tab 0 0

 unit _ = Cols 1

Example: Parsing UTF8

- Valid UTF8 encoded Chars have the form:
 - [0x00...0x7F]
 - [0xC0...0xDF] extra
 - [0xE0...0xEF] extra extra
 - [0xF0...0xF₄] extra extra extra
- where extra = [0x80...0xBF] contains 6 bits of info in the LSBs and the only valid representation is the shortest one for each symbol.

UTF8 as a Reducer Transformer

data UTF8 m = Segment !Prefix m !Suffix | Chunk !Suffix

instance (Char `Reducer` m) => Monoid (UTF8 m)

where ...

instance (Char `Reducer` m) => (Byte `Reducer` UTF8 m)

where ...

Given 7 bytes we must have seen a full Char.

We only need track up to 3 bytes on either side.

Putting the pieces together so far

We can:

- Parse a file as a Lazy ByteString,
- Ignore alignment of the chunks and parse UTF8, automatically cleaning up the ends as needed when we glue the reductions of our chunks together.
- We can feed that into a complicated Char `Reducer` that uses modular components like Delta.

Compressive Parsing

- LZ78 decompression never compares values in the dictionary. Decompress **in** the monoid, caching the results.
- Unlike later refinements (LZW, LZSS, etc.) LZ78 doesn't require every value to initialize the dictionary permitting infinite alphabets (i.e. Integers)
- We can compress chunkwise, permitting parallelism
- Decompression fits on a slide.

Compressive Parsing

```
newtype LZ78 a = LZ78 [Token a]
```

```
data Token a = Token a !Int
```

```
instance Generator (LZ78 a) where
```

```
  type Elem (LZ78 a) = a
```

```
  mapTo f m (LZ78 xs) = mapTo' f m (Seq.singleton mempty) xs
```

```
mapTo' :: (e `Reducer` m) => (a -> e) -> m -> Seq m -> [Token a] -> m
```

```
mapTo' _ m _ [] = m
```

```
mapTo' f m s (Token c w:ws) = m `mappend` mapTo' f v (s |> v) ws
```

```
  where v = Seq.index s w `snoc` f c
```

Other Compressive Parsers

- The dictionary size in the previous example can be bounded, so we can provide reuse of common monoids **up to** a given size or within a given window.
- Other extensions to LZW (i.e. LZAP) can be adapted to LZ78, and work even better over monoids than normal!
- Bentley-McIlroy (the basis of bmdiff and open-vcdiff) can be used to reuse all common submonoids **over** a given size.

Going Deeper

Algebraic Structure Provides Opportunity

Structure	Example Opportunity
Semigroup	Parallelized Folds
Monoid	Unit
Group	Inverses/Undo
Commutative Monoid	Reordering Computation
Applicative	Synthesized Attributes
Abelian Group	Out-Of-Order Undo
Ringoid	Cancellative Zero
Right Seminearring	Context-Free Recognizers
Alternative	Context-Free Attribute Grammars
Monad	Context-Sensitivity



Conclusion

- Monoids are *everywhere*
 - Reducers allow *efficient* use of Monoids
 - Generators can apply Reducers in *parallel*
 - Monoids/Reducers are *composable*
 - Compression can improve *performance*
 - Algebraic structures provide *opportunity*
- 